

# Linear Programming: Linear Programming

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Advanced Algorithms and Complexity  
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# Learning Objectives

- Understand the formal definition of a linear programming problem.
- Provide some examples of linear programming problems

# Last Time

Factory. Set  $M, W$  to maximize  $200M + 100W$  subject to

- $W \geq 0$ .
- $100 \geq M \geq 0$ .
- $W \geq 2M$ .
- $100,000 \geq 200(W - 2M) + 600M$ .

# Linear Programming

Linear programming asks for real numbers  $x_1, x_2, \dots, x_n$  satisfying linear inequalities:

$$a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n \geq b_1$$

...

$$a_{m1}x_1 + a_{m2}x_2 + \dots + a_{mn}x_n \geq b_m$$

So that a linear objective

$$v_1x_1 + v_2x_2 + \dots + v_nx_n$$

is as large (or small) as possible.

# Notation

## Linear Programming

**Input:** An  $m \times n$  matrix  $A$  and vectors  $b \in \mathbb{R}^m, v \in \mathbb{R}^n$

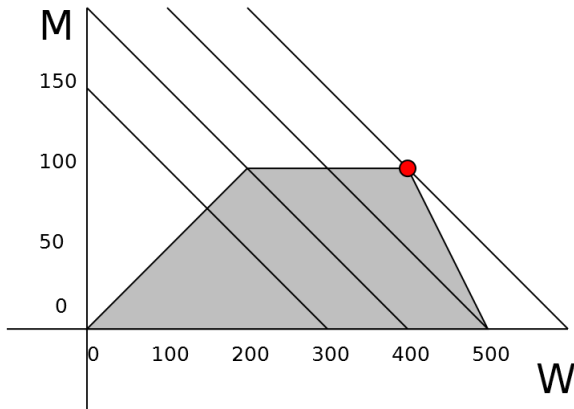
**Output:** A vector  $x \in \mathbb{R}^n$  so that  $Ax \geq b$  and  $v \cdot x$  is as large (or small) as possible.

# Examples

Linear programming is useful because an extraordinary number of problems can be put into this framework.

# Factory Example

The factory example we just worked.



# The Diet Problem

Studied by George Stigler in the 1930s and 1940s.

How cheaply can you purchase food for a healthy diet?



# Variables

You have a number of types of food (bread, milk, apples, etc.).

For each you have a variable giving the number of servings per day.

$$x_{bread}, x_{milk}, x_{apples}, \dots$$

# Constraints

Non-negative number of servings:

$$x_f \geq 0.$$

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Sufficient calories/day:

$$\begin{aligned} & (\text{Cal/serving bread})x_{\text{bread}} \\ & + (\text{Cal/serving milk})x_{\text{milk}} + \dots \geq 2000. \end{aligned}$$

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Similar constraints for other nutritional needs  
(vitamin C, protein, etc.)

# Optimization

Minimize cost.

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Warning: actually doing this can get you some pretty weird diets.

# Network Flow

Network flow problems are actually just a **special case** of linear programming problems!

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Objective:

$$\sum_{e \text{ out of } s} f_e - \sum_{e \text{ into } s} f_e$$

# Strange Cases

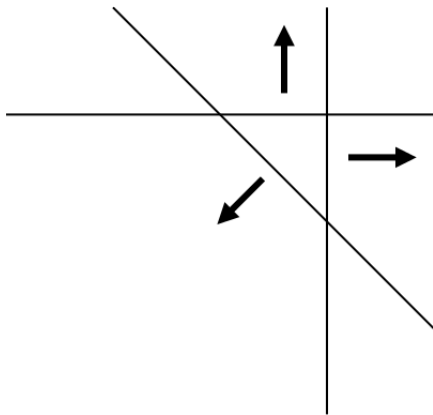
There are a couple of edge cases to keep in mind here.

- No Solution
- No Optimum

# No Solution

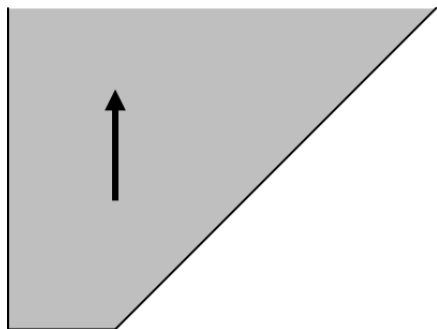
Consider the system:

$$x \geq 1, \quad y \geq 1, \quad x + y \leq 1.$$



# No Optimum

Consider trying to maximize  $x$  subject to  $x \geq 0$ ,  $y \geq 0$ , and  $x - y \geq 1$ .



# Problem

Of these three systems, one has no solution, one has no maximum  $x$  value, and one has a maximum. Which is which?

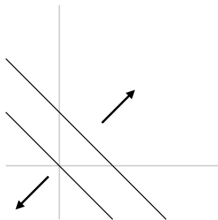
(A)  $x + y \geq 1, x + y \leq 0.$

(B)  $x + y \leq 2, x - y \leq 1.$

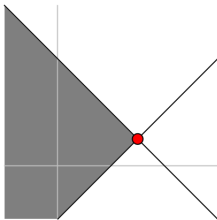
(C)  $x + y \geq 0, x - y \leq 0.$

# Solution

A



B



C

